




# Listing Prime Numbers Periodically

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## Abstract

Over millennia, people have considered prime numbers as unpredictable and hard to list. This study confirms that all positive integers without factors of 2, 3, 5 and 7 can be regarded as the offspring of 48 roots in the interval [11, 220]. These roots are to generate a Factor-Pair Table composed of 48 columns, each of which forms a basic binary matrix to indicate the locations of composite numbers. For a given interval, these composites are removed via mapping effect. After primality checking of the rest of the numbers, we identify all primes within the interval. We then form a Formula of Primes for prime prediction. Moreover, the Periodic Table of Primes (PTP) [1] is reconstructed based on the proposed algorithm to illustrate the distribution of primes and composites.

## Keywords

Prime Listing, Factor-Pair Table, Mapping Effect

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## 1. Introduction

The problem of listing all primes up to a given limit has always been a challenging task. The prime number theorem [2] states that the prime-counting function  $\pi(n)$  is approximated by  $n/\log n$  when  $n$  is large, but listing all primes in a given interval has been considered as a much more difficult problem.

Except for trial division, which is too time-consuming, most of the existing algorithms are designed for finding all primes from the start, namely, from 2 to the given limit. The ancient Sieve of Eratosthenes iteratively removes composites as the multiples of primes already found [3]. In one of his notes [4], Dijkstra proposed a prime number algorithm that is more space-saving than the Sieve of Eratosthenes and has lower time complexity than trial division. Among other variants

of the Sieve of Eratosthenes, Atkin's [5] and Pritchard's [6] improve the algorithm with preprocessing efforts, but require a trade-off between time and space complexity.

In contrast, little has been done on the prime listing within a given interval that starts from a large number. Helfgott [7] introduced an improved version of the Sieve of Eratosthenes that can construct all primes in an interval  $[n - \Delta, n + \Delta]$  in space  $O(\Delta)$  and time  $O(\Delta \log n)$ .

In this paper, we propose an efficient algorithm to list primes in a given interval that starts from any number. The algorithm is based on the Periodic Table of Primes (PTP) developed by Li *et al.* [1]. Although similar roots were proposed in 1982 as Wheel factorization [8], the PTP [1] is an infinite table generated from the Cyclic Table of Composites through rigorous mathematical proofs. Based on the algorithm presented in this paper, the PTP [1] is further reconstructed.

Our algorithm first restricts the numbers to be tested to a smaller subset, specifically those in the given interval and the PTP. Then, the mapping effect determines that, in expectation, more than 54.75% of them are composite numbers in constant time. Combined with existing efficient primality test algorithms, including the AKS primality test [9] and Miller-Rabin primality test [10] [11], it reveals the primality of the rest of the numbers and lists all primes in that given interval.

## Notations

- $N_0 = \{0, 1, 2, \dots\}$ .
- $N_+ = \{1, 2, 3, \dots\}$ .
- $\Lambda = \{n \in N_+ \mid n > 1 \text{ and } n \text{ has no factors of } 2, 3, 5 \text{ and } 7\}$ .
- $S = \{r_i \mid r_i \text{ is an integer, } 11 \leq r_i \leq 220, r_i \in \Lambda\}$   
 $= \{r_1 = 11, r_2 = 13, r_3 = 17, \dots, r_{48} = 211\}$ .
- $Q = \{q_j \mid q_j = r_j \in S\} = \{q_1 = 11, q_2 = 13, q_3 = 17, \dots, q_{48} = 211\}$ .
- $[a, b]_N = \{n \in N \mid a \leq n \leq b\}$ .
- $\lceil x \rceil$  = ceiling function of the smallest integer greater than or equal to  $x$ .
- $\lfloor x \rfloor$  = floor function of the largest integer less than  $x$ .

## 2. Basic Theory

This study proposes a novel way, called Periodic-Table-of-Primes (PTP) method, to list all primes within a pre-specified integer interval. The deduction of this method is based on the following observations:

### 2.1. Roots of Integers

By choosing the first four primes, 2, 3, 5 and 7, we can group every  $2 \times 3 \times 5 \times 7 = 210$  numbers in one period starting from 11, *i.e.*,  $[11, 220]_N, [221, 430]_N, [431, 640]_N, \dots$ . Within each of these periods, there are 48 numbers without factors of 2, 3, 5 and 7. In particular, in  $[11, 220]_N$ , there exist 48 numbers called roots, and we denote

$$S = \{r_1 = 11, r_2 = 13, r_3 = 17, r_4 = 19, r_5 = 23, r_6 = 29, r_7 = 31, r_8 = 37, r_9 = 41, r_{10} = 43, r_{11} = 47, r_{12} = 53, r_{13} = 59, r_{14} = 61, r_{15} = 67, r_{16} = 71, r_{17} = 73, r_{18} = 79, r_{19} = 83, r_{20} = 89, r_{21} = 97, r_{22} = 101, r_{23} = 103, r_{24} = 107, r_{25} = 109, r_{26} = 113, r_{27} = 121, r_{28} = 127, r_{29} = 131, r_{30} = 137, r_{31} = 139, r_{32} = 143, r_{33} = 149, r_{34} = 151, r_{35} = 157, r_{36} = 163, r_{37} = 167, r_{38} = 169, r_{39} = 173, r_{40} = 179, r_{41} = 181, r_{42} = 187, r_{43} = 191, r_{44} = 193, r_{45} = 197, r_{46} = 199, r_{47} = 209, r_{48} = 211\}.$$

**Theorem 1.** For any number  $\alpha \in \Lambda$ , there exists a unique  $r_i \in S$  and  $k \in N_0$  such that

$$\alpha = r_i + 210k. \tag{1}$$

From Theorem 1, we know that each number  $\alpha \in \Lambda$  has a root and a period. In other words,  $\alpha$  is the  $k^{\text{th}}$  descendant of  $r_i$ .

From notations indicating that  $Q = S$  and  $q_j = r_j, j = 1, 2, \dots, 48$ , we can also know that any factor of  $\alpha$  can be expressed as  $q_j + 210\theta$ , for some  $q_j \in Q$  and  $\theta \in N_0$ . The following observation further analyzes the structure of  $\alpha$ 's factors.

### 2.2. Factor-Pair Table with Generation Information

Given  $q_j \in Q$  and  $r_i \in S$ , denote  $q_{j|i}$  as the unique element in  $Q$  such that  $q_j \times q_{j|i} - r_i$  is a multiplier of 210, where  $i, j \in \{1, 2, \dots, 48\}$ . Here  $(q_j, q_{j|i})$  is called a factor-pair with respect to  $r_i$ . We also observe there exists an integer

$$t_{j|i} = \frac{q_j q_{j|i} - r_i}{210} \tag{2}$$

to specify the generation of  $r_i$  in which the descendant of  $q_j$  and  $q_{j|i}$ , i.e.,  $q_j \times q_{j|i}$  lies.

Collecting all  $(q_j, q_{j|i})$  and  $t_{j|i}$  for all  $i, j \in \{1, 2, \dots, 48\}$  together, we then form a Factor-Pair Table with generation information. It is a  $48 \times 48$  table made of element  $(q_j, q_{j|i}), t_{j|i}$ .

We denote each column of the table as  $V_j$ , where

$$V_j = \left( (q_j, q_{j|1}), t_{j|1}; \dots; (q_j, q_{j|i}), t_{j|i}; \dots; (q_j, q_{j|48}), t_{j|48} \right)^T \text{ for } j = 1, 2, 3, \dots, 48.$$

Take  $j = 1$  for instance,

$$V_1 = \left( (11, 211), 11; (11, 173), 9; (11, 97), 5; (11, 59), 3; \dots; (11, 191), 9 \right)^T,$$

where  $(q_1, q_{1|1}, t_{1|1}) = (q_1, q_{48}, t_{1|1}) = (11, 211, 11)$ , as shown in **Table 1**.

The Factor-Pair Table with generation information has the following properties:

- 1) In each row of the table, there are 48 factor-pairs.
- 2) If  $q_j = q_{j|i}$ , then  $(q_j, q_{j|i})$  is called a co-factor-pair. Co-factor-pairs only exist in 6 rows of the table. In each row corresponding to  $r_{18} = 79, r_{25} = 109, r_{27} = 121, r_{34} = 151, r_{38} = 169$  and  $r_{48} = 211$ , there are 8 co-factor-pairs among the 48 factor-pairs.

**Table 1.** Factor-pair table with generation information.

$i$	$r_i$	$q_j$															
		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16
1	11	(11, 211), 11	(13, 17), 1	(17, 13), 1	(19, 89), 8	(23, 37), 4	(29, 109), 15	(31, 41), 6	(37, 23), 4	(41, 31), 6	(43, 137), 28	(47, 103), 23	(53, 127), 32	(59, 139), 39	(61, 131), 38	(67, 113), 36	(71, 151), 51
2	13	(11, 173), 9	(13, 211), 13	(17, 149), 12	(19, 67), 6	(23, 101), 11	(29, 167), 23	(31, 163), 24	(37, 199), 35	(41, 113), 22	(43, 181), 37	(47, 179), 40	(53, 131), 33	(59, 107), 30	(61, 193), 56	(67, 19), 6	(71, 83), 28
3	17	(11, 97), 5	(13, 179), 11	(17, 211), 17	(19, 23), 2	(23, 19), 2	(29, 73), 10	(31, 197), 29	(37, 131), 23	(41, 67), 13	(43, 59), 12	(47, 121), 27	(53, 139), 35	(59, 43), 12	(61, 107), 31	(67, 41), 13	(71, 157), 53
4	19	(11, 59), 3	(13, 163), 10	(17, 137), 11	(19, 211), 19	(23, 83), 9	(29, 131), 18	(31, 109), 16	(37, 97), 17	(41, 149), 29	(43, 103), 21	(47, 197), 44	(53, 143), 36	(59, 11), 3	(61, 169), 49	(67, 157), 50	(71, 89), 30
5	23	(11, 193), 10	(13, 131), 8	(17, 199), 16	(19, 167), 15	(23, 211), 23	(29, 37), 5	(31, 143), 21	(37, 29), 5	(41, 103), 20	(43, 191), 39	(47, 139), 31	(53, 151), 38	(59, 157), 44	(61, 83), 24	(67, 179), 57	(71, 163), 55
6	29	(11, 79), 4	(13, 83), 5	(17, 187), 15	(19, 101), 9	(23, 193), 21	(29, 211), 29	(31, 89), 13	(37, 137), 24	(41, 139), 27	(43, 113), 23	(47, 157), 35	(53, 163), 41	(59, 61), 17	(61, 59), 17	(67, 107), 34	(71, 169), 57
7	31	(11, 41), 2	(13, 67), 4	(17, 113), 9	(19, 79), 7	(23, 47), 5	(29, 59), 8	(31, 211), 31	(37, 103), 18	(41, 11), 2	(43, 157), 32	(47, 23), 5	(53, 167), 42	(59, 29), 8	(61, 121), 35	(67, 13), 4	(71, 101), 34
8	37	(11, 137), 7	(13, 19), 1	(17, 101), 8	(19, 13), 1	(23, 29), 3	(29, 23), 3	(31, 157), 23	(37, 211), 37	(41, 47), 9	(43, 79), 16	(47, 41), 9	(53, 179), 45	(59, 143), 40	(61, 97), 28	(67, 151), 48	(71, 107), 36
9	41	(11, 61), 3	(13, 197), 12	(17, 163), 13	(19, 179), 16	(23, 157), 17	(29, 139), 19	(31, 191), 28	(37, 143), 25	(41, 211), 41	(43, 167), 34	(47, 193), 43	(53, 187), 47	(59, 79), 22	(61, 11), 3	(67, 173), 55	(71, 181), 61
10	43	(11, 23), 1	(13, 181), 11	(17, 89), 7	(19, 157), 14	(23, 11), 1	(29, 197), 27	(31, 103), 15	(37, 109), 19	(41, 83), 16	(43, 211), 43	(47, 59), 13	(53, 191), 48	(59, 47), 13	(61, 73), 21	(67, 79), 25	(71, 113), 38
11	47	(11, 157), 8	(13, 149), 9	(17, 151), 12	(19, 113), 10	(23, 139), 15	(29, 103), 14	(31, 137), 20	(37, 41), 7	(41, 37), 7	(43, 89), 18	(47, 211), 47	(53, 199), 50	(59, 193), 54	(61, 197), 57	(67, 101), 32	(71, 187), 63
12	53	(11, 43), 2	(13, 101), 6	(17, 139), 11	(19, 47), 4	(23, 121), 13	(29, 67), 9	(31, 83), 12	(37, 149), 26	(41, 73), 14	(43, 11), 2	(47, 19), 4	(53, 211), 53	(59, 97), 27	(61, 173), 50	(67, 29), 9	(71, 193), 65
13	59	(11, 139), 7	(13, 53), 3	(17, 127), 10	(19, 191), 17	(23, 103), 11	(29, 31), 4	(31, 29), 4	(37, 47), 8	(41, 109), 21	(43, 143), 29	(47, 37), 8	(53, 13), 3	(59, 211), 59	(61, 149), 43	(67, 167), 53	(71, 199), 67
14	61	(11, 101), 5	(13, 37), 2	(17, 53), 4	(19, 169), 15	(23, 167), 18	(29, 89), 12	(31, 151), 22	(37, 13), 2	(41, 191), 37	(43, 187), 38	(47, 113), 25	(53, 17), 4	(59, 179), 50	(61, 211), 61	(67, 73), 23	(71, 131), 44
15	67	(11, 197), 10	(13, 199), 12	(17, 41), 3	(19, 103), 9	(23, 149), 16	(29, 53), 7	(31, 97), 14	(37, 121), 21	(41, 17), 3	(43, 109), 22	(47, 131), 29	(53, 29), 7	(59, 83), 23	(61, 187), 54	(67, 211), 67	(71, 137), 46
16	71	(11, 121), 6	(13, 167), 10	(17, 103), 8	(19, 59), 5	(23, 67), 7	(29, 169), 23	(31, 131), 19	(37, 53), 9	(41, 181), 35	(43, 197), 40	(47, 73), 16	(53, 37), 9	(59, 19), 5	(61, 101), 29	(67, 23), 7	(71, 211), 71
17	73	(11, 83), 4	(13, 151), 9	(17, 29), 2	(19, 37), 3	(23, 131), 14	(29, 17), 2	(31, 43), 6	(37, 19), 3	(41, 53), 10	(43, 31), 6	(47, 149), 33	(53, 41), 10	(59, 197), 55	(61, 163), 47	(67, 139), 44	(71, 143), 48
18	79	(11, 179), 9	(13, 103), 6	(17, 17), 1	(19, 181), 16	(23, 113), 12	(29, 191), 26	(31, 199), 29	(37, 127), 22	(41, 89), 17	(43, 163), 33	(47, 167), 37	(53, 53), 13	(59, 101), 28	(61, 139), 40	(67, 67), 21	(71, 149), 50
19	83	(11, 103), 5	(13, 71), 4	(17, 79), 6	(19, 137), 12	(23, 31), 3	(29, 97), 13	(31, 23), 3	(37, 59), 10	(41, 43), 8	(43, 41), 8	(47, 109), 24	(53, 61), 15	(59, 37), 10	(61, 53), 15	(67, 89), 28	(71, 13), 4
20	89	(11, 199), 10	(13, 23), 1	(17, 67), 5	(19, 71), 6	(23, 13), 1	(29, 61), 8	(31, 179), 26	(37, 167), 29	(41, 79), 15	(43, 173), 35	(47, 127), 28	(53, 73), 18	(59, 151), 42	(61, 29), 8	(67, 17), 5	(71, 19), 6
21	97	(11, 47), 2	(13, 169), 10	(17, 191), 15	(19, 193), 17	(23, 59), 6	(29, 83), 11	(31, 37), 5	(37, 31), 5	(41, 197), 38	(43, 139), 28	(47, 11), 2	(53, 89), 22	(59, 23), 6	(61, 67), 19	(67, 61), 19	(71, 167), 56
22	101	(11, 181), 9	(13, 137), 8	(17, 43), 3	(19, 149), 13	(23, 187), 20	(29, 199), 27	(31, 71), 10	(37, 173), 30	(41, 151), 29	(43, 17), 3	(47, 163), 36	(53, 97), 24	(59, 169), 47	(61, 191), 55	(67, 83), 26	(71, 31), 10
23	103	(11, 143), 7	(13, 121), 7	(17, 179), 14	(19, 127), 11	(23, 41), 4	(29, 47), 6	(31, 193), 28	(37, 139), 24	(41, 23), 4	(43, 61), 12	(47, 29), 6	(53, 101), 25	(59, 137), 38	(61, 43), 12	(67, 199), 63	(71, 173), 58
24	107	(11, 67), 3	(13, 89), 5	(17, 31), 2	(19, 83), 7	(23, 169), 18	(29, 163), 22	(31, 17), 2	(37, 71), 12	(41, 187), 36	(43, 149), 30	(47, 181), 40	(53, 109), 27	(59, 73), 20	(61, 167), 48	(67, 11), 3	(71, 37), 12
25	109	(11, 29), 1	(13, 73), 4	(17, 167), 13	(19, 61), 5	(23, 23), 2	(29, 11), 1	(31, 139), 20	(37, 37), 6	(41, 59), 11	(43, 193), 39	(47, 47), 10	(53, 113), 28	(59, 41), 11	(61, 19), 5	(67, 127), 40	(71, 179), 60
26	113	(11, 163), 8	(13, 41), 2	(17, 19), 1	(19, 17), 1	(23, 151), 16	(29, 127), 17	(31, 173), 25	(37, 179), 31	(41, 13), 2	(43, 71), 14	(47, 199), 44	(53, 121), 30	(59, 187), 52	(61, 143), 41	(67, 149), 47	(71, 43), 14
27	121	(11, 11), 0	(13, 187), 11	(17, 143), 11	(19, 139), 12	(23, 197), 21	(29, 149), 20	(31, 31), 4	(37, 43), 7	(41, 131), 25	(43, 37), 7	(47, 83), 18	(53, 137), 34	(59, 59), 16	(61, 181), 52	(67, 193), 61	(71, 191), 64
28	127	(11, 107), 5	(13, 139), 8	(17, 131), 10	(19, 73), 6	(23, 179), 19	(29, 113), 15	(31, 187), 27	(37, 151), 26	(41, 167), 32	(43, 169), 34	(47, 101), 22	(53, 149), 37	(59, 173), 48	(61, 157), 45	(67, 121), 38	(71, 197), 66
29	131	(11, 31), 1	(13, 107), 6	(17, 193), 15	(19, 29), 2	(23, 97), 10	(29, 19), 2	(31, 11), 1	(37, 83), 14	(41, 121), 23	(43, 47), 9	(47, 43), 9	(53, 157), 39	(59, 109), 30	(61, 71), 20	(67, 143), 45	(71, 61), 20
30	137	(11, 127), 6	(13, 59), 3	(17, 181), 14	(19, 173), 15	(23, 79), 8	(29, 193), 26	(31, 167), 24	(37, 191), 33	(41, 157), 30	(43, 179), 36	(47, 61), 13	(53, 169), 42	(59, 13), 3	(61, 47), 13	(67, 71), 22	(71, 67), 22
31	139	(11, 89), 4	(13, 43), 2	(17, 107), 8	(19, 151), 13	(23, 143), 15	(29, 41), 5	(31, 79), 11	(37, 157), 27	(41, 29), 5	(43, 13), 2	(47, 137), 30	(53, 173), 43	(59, 191), 53	(61, 109), 31	(67, 187), 59	(71, 209), 70
32	143	(11, 13), 0	(13, 11), 0	(17, 169), 13	(19, 107), 9	(23, 61), 6	(29, 157), 21	(31, 113), 16	(37, 89), 15	(41, 193), 37	(43, 101), 20	(47, 79), 17	(53, 181), 45	(59, 127), 35	(61, 23), 6	(67, 209), 66	(71, 73), 24
33	149	(11, 109), 5	(13, 173), 10	(17, 157), 12	(19, 41), 3	(23, 43), 4	(29, 121), 16	(31, 59), 8	(37, 197), 34	(41, 19), 3	(43, 23), 4	(47, 97), 21	(53, 193), 48	(59, 31), 8	(61, 209), 60	(67, 137), 43	(71, 79), 26
34	151	(11, 71), 3	(13, 157), 9	(17, 83), 6	(19, 19), 1	(23, 107), 11	(29, 179), 24	(31, 181), 26	(37, 163), 28	(41, 101), 19	(43, 67), 13	(47, 173), 38	(53, 197), 49	(59, 209), 58	(61, 61), 17	(67, 43), 13	(71, 11), 3
35	157	(11, 167), 8	(13, 109), 6	(17, 71), 5	(19, 163), 14	(23, 89), 9	(29, 143), 19	(31, 127), 18	(37, 61), 10	(41, 137), 26	(43, 199), 40	(47, 191), 42	(53, 209), 52	(59, 113), 31	(61, 37), 10	(67, 181), 57	(71, 17), 5
36	163	(11, 53), 2	(13, 61), 3	(17, 59), 4	(19, 97), 8	(23, 71), 7	(29, 107), 14	(31, 73), 10	(37, 169), 29	(41, 173), 33	(43, 121), 24	(47, 209), 46	(53, 11), 2	(59, 17), 4	(61, 13), 3	(67, 109), 34	(71, 23), 7
37	167	(11, 187), 9	(13, 29), 1	(17, 121), 9	(19, 53), 4	(23, 199), 21	(29, 13), 1	(31, 107), 15	(37, 101), 17	(41, 127), 24	(43, 209), 42	(47, 151), 33	(53, 19), 4	(59, 163), 45	(61, 137), 39	(67, 131), 41	(71, 97), 32
38	169	(11, 149), 7	(13, 13), 0	(17, 47), 3	(19, 31), 2	(23, 53), 5	(29, 71), 9	(31, 19), 2	(37, 67), 11	(41, 209), 40	(43, 43), 8	(47, 17), 3	(53, 23), 5	(59, 131), 36	(61, 199), 57	(67, 37), 11	(71, 29), 9
39	173	(11, 73), 3	(13, 191), 11	(17, 109), 8	(19, 197), 17	(23, 181), 19	(29, 187), 25	(31, 53), 7	(37, 209), 36	(41, 163), 31	(43, 131), 26	(47, 169), 37	(53, 31), 7	(59, 67), 18	(61, 113), 32	(67, 59), 18	(71, 103), 34
40	179	(11, 169), 8	(13, 143), 8	(17, 97), 7	(19, 131), 11	(23, 163), 17	(29, 151), 20	(31, 209), 30	(37, 107), 18	(41, 199), 38	(43, 53), 10	(47, 187), 41	(53, 43), 10	(59, 181), 50	(61, 89), 25	(67, 197), 62	(71, 109), 36
41	181	(11, 131), 6	(13, 127), 7	(17, 23), 1	(19, 109), 9	(23, 17), 1	(29, 209), 28	(31, 121), 17	(37, 73), 12	(41, 71), 13	(43, 97), 19	(47, 53), 11	(53, 47), 11	(59, 149), 41	(61, 151), 43	(67, 103), 32	(71, 41), 13
42	187	(11, 17), 0	(13, 79), 4	(17, 11), 0	(19, 43), 3	(23, 209), 22	(29, 173), 23	(31, 67), 9	(37, 181), 31	(41, 107), 20	(43, 19), 3	(47, 71), 15	(53, 59), 14	(59, 53), 14	(61, 127), 36	(67, 31), 9	(71, 47), 15
43	191	(11, 151), 7	(13, 47), 2	(17, 73), 5	(19, 209), 18	(23, 127), 13	(29, 79), 10	(31, 101), 14	(37, 113), 19	(41, 61), 11	(43, 107), 21	(47					

Continued

		j																														
		17	18	19	20	21	22	23	24	25	26	27	28	29	30	31	32															
i	r <sub>i</sub>	q <sub>j</sub>																														
		73	79	83	89	97	101	103	107	109	113	121	127	131	137	139	143															
1	11	(73, 167), 58	(79, 149), 56	(83, 157), 62	(89, 149), 8	(97, 143), 66	(101, 181), 87	(103, 47), 23	(107, 163), 83	(109, 29), 15	(113, 67), 36	(121, 191), 110	(127, 53), 32	(131, 61), 38	(137, 43), 28	(139, 59), 39	(143, 97), 66															
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3	17	(73, 29), 10	(79, 173), 65	(83, 109), 43	(89, 163), 69	(97, 11), 5	(101, 127), 61	(103, 149), 73	(107, 61), 31	(109, 83), 43	(113, 199), 107	(121, 47), 27	(127, 101), 61	(131, 37), 23	(137, 181), 118	(139, 53), 35	(143, 169), 115															
4	19	(73, 193), 67	(79, 181), 68	(83, 23), 9	(89, 71), 30	(97, 37), 17	(101, 179), 86	(103, 43), 21	(107, 167), 85	(109, 31), 16	(113, 173), 93	(121, 139), 80	(127, 187), 113	(131, 29), 18	(137, 17), 11	(139, 121), 80	(143, 53), 36															
5	23	(73, 101), 35	(79, 197), 74	(83, 61), 24	(89, 97), 41	(97, 89), 41	(101, 73), 35	(103, 41), 20	(107, 169), 86	(109, 137), 71	(113, 121), 65	(121, 113), 65	(127, 149), 90	(131, 13), 8	(137, 109), 71	(139, 47), 31	(143, 31), 21															
6	29	(73, 173), 60	(79, 11), 4	(83, 13), 5	(89, 31), 13	(97, 167), 77	(101, 19), 9	(103, 143), 70	(107, 67), 34	(109, 191), 99	(113, 43), 23	(121, 179), 103	(127, 197), 119	(131, 199), 124	(137, 37), 24	(139, 41), 27	(143, 103), 70															
7	31	(73, 127), 44	(79, 19), 7	(83, 137), 54	(89, 149), 63	(97, 193), 89	(101, 71), 34	(103, 37), 18	(107, 173), 88	(109, 139), 72	(113, 17), 9	(121, 61), 35	(127, 73), 44	(131, 191), 119	(137, 83), 54	(139, 109), 72	(143, 197), 134															
8	37	(73, 199), 69	(79, 43), 16	(83, 89), 35	(89, 83), 35	(97, 61), 28	(101, 17), 8	(103, 139), 68	(107, 71), 36	(109, 193), 100	(113, 149), 80	(121, 127), 73	(127, 121), 73	(131, 167), 104	(137, 11), 7	(139, 103), 68	(143, 59), 40															
9	41	(73, 107), 37	(79, 59), 22	(83, 127), 50	(89, 109), 46	(97, 113), 52	(101, 121), 58	(103, 137), 67	(107, 73), 37	(109, 89), 46	(113, 97), 52	(121, 101), 58	(127, 83), 50	(131, 151), 94	(137, 103), 67	(139, 29), 19	(143, 37), 25															
10	43	(73, 61), 21	(79, 67), 25	(83, 41), 16	(89, 17), 7	(97, 139), 64	(101, 173), 83	(103, 31), 15	(107, 179), 91	(109, 37), 19	(113, 71), 38	(121, 193), 111	(127, 169), 102	(131, 143), 89	(137, 149), 97	(139, 97), 64	(143, 131), 89															
11	47	(73, 179), 62	(79, 83), 31	(83, 79), 31	(89, 43), 18	(97, 191), 88	(101, 67), 32	(103, 29), 14	(107, 181), 92	(109, 143), 74	(113, 19), 10	(121, 167), 96	(127, 131), 79	(131, 127), 79	(137, 31), 20	(139, 23), 15	(143, 109), 74															
12	53	(73, 41), 14	(79, 107), 40	(83, 31), 12	(89, 187), 79	(97, 59), 27	(101, 13), 6	(103, 131), 64	(107, 79), 40	(109, 197), 102	(113, 151), 81	(121, 23), 13	(127, 179), 108	(131, 103), 64	(137, 169), 110	(139, 17), 11	(143, 181), 123															
13	59	(73, 113), 39	(79, 131), 49	(83, 193), 76	(89, 121), 51	(97, 137), 63	(101, 169), 81	(103, 23), 11	(107, 187), 95	(109, 41), 21	(113, 73), 39	(121, 89), 51	(127, 17), 10	(131, 79), 49	(137, 97), 63	(139, 11), 7	(143, 43), 29															
14	61	(73, 67), 23	(79, 139), 52	(83, 107), 42	(89, 29), 12	(97, 163), 75	(101, 11), 5	(103, 127), 62	(107, 83), 42	(109, 199), 103	(113, 47), 25	(121, 181), 104	(127, 103), 62	(131, 71), 44	(137, 143), 93	(139, 79), 52	(143, 137), 93															
15	67	(73, 139), 48	(79, 163), 61	(83, 59), 23	(89, 173), 73	(97, 31), 14	(101, 167), 80	(103, 19), 9	(107, 191), 97	(109, 43), 22	(113, 179), 96	(121, 37), 21	(127, 151), 91	(131, 47), 29	(137, 71), 46	(139, 73), 48	(143, 209), 142															
16	71	(73, 47), 16	(79, 179), 67	(83, 97), 38	(89, 199), 84	(97, 83), 38	(101, 61), 29	(103, 17), 8	(107, 193), 98	(109, 149), 77	(113, 127), 68	(121, 11), 6	(127, 113), 68	(131, 31), 19	(137, 163), 106	(139, 209), 138	(143, 187), 127															
17	73	(73, 211), 73	(79, 187), 70	(83, 11), 4	(89, 107), 45	(97, 109), 50	(101, 113), 54	(103, 121), 59	(107, 89), 45	(109, 97), 50	(113, 101), 54	(121, 103), 59	(127, 199), 120	(131, 23), 14	(137, 209), 136	(139, 67), 44	(143, 71), 48															
18	79	(73, 73), 25	(79, 211), 79	(83, 173), 68	(89, 41), 17	(97, 187), 86	(101, 59), 28	(103, 13), 6	(107, 197), 100	(109, 151), 78	(113, 23), 12	(121, 169), 97	(127, 37), 22	(131, 209), 130	(137, 137), 89	(139, 61), 40	(143, 143), 97															
19	83	(73, 191), 66	(79, 17), 6	(83, 211), 83	(89, 67), 28	(97, 29), 13	(101, 163), 78	(103, 11), 5	(107, 199), 101	(109, 47), 24	(113, 181), 97	(121, 143), 82	(127, 209), 126	(131, 193), 120	(137, 19), 12	(139, 197), 130	(143, 121), 82															
20	89	(73, 53), 18	(79, 41), 15	(83, 163), 64	(89, 211), 89	(97, 107), 49	(101, 109), 52	(103, 113), 55	(107, 97), 49	(109, 101), 52	(113, 103), 55	(121, 209), 120	(127, 47), 28	(131, 169), 105	(137, 157), 102	(139, 191), 126	(143, 193), 131															
21	97	(73, 79), 27	(79, 73), 27	(83, 29), 11	(89, 53), 22	(97, 211), 97	(101, 107), 51	(103, 109), 53	(107, 101), 51	(109, 103), 53	(113, 209), 112	(121, 157), 90	(127, 181), 109	(131, 137), 85	(137, 131), 85	(139, 43), 28	(143, 149), 101															
22	101	(73, 197), 68	(79, 89), 33	(83, 67), 26	(89, 79), 33	(97, 53), 24	(101, 211), 101	(103, 107), 52	(107, 103), 52	(109, 209), 108	(113, 157), 84	(121, 131), 75	(127, 143), 86	(131, 121), 75	(137, 13), 8	(139, 179), 118	(143, 127), 86															
23	103	(73, 151), 52	(79, 97), 36	(83, 191), 75	(89, 197), 83	(97, 79), 36	(101, 53), 25	(103, 211), 103	(107, 209), 106	(109, 157), 81	(113, 131), 70	(121, 13), 7	(127, 19), 11	(131, 113), 70	(137, 59), 38	(139, 37), 24	(143, 11), 7															
24	107	(73, 59), 20	(79, 113), 42	(83, 19), 7	(89, 13), 5	(97, 131), 60	(101, 157), 75	(103, 209), 102	(107, 211), 107	(109, 53), 27	(113, 79), 42	(121, 197), 113	(127, 191), 115	(131, 97), 60	(137, 151), 98	(139, 173), 114	(143, 199), 135															
25	109	(73, 13), 4	(79, 121), 45	(83, 143), 56	(89, 131), 55	(97, 157), 72	(101, 209), 100	(103, 103), 50	(107, 107), 54	(109, 211), 109	(113, 53), 28	(121, 79), 45	(127, 67), 40	(131, 89), 55	(137, 197), 128	(139, 31), 20	(143, 83), 56															
26	113	(73, 131), 45	(79, 137), 51	(83, 181), 71	(89, 157), 66	(97, 209), 96	(101, 103), 49	(103, 101), 49	(107, 109), 55	(109, 107), 55	(113, 211), 113	(121, 53), 30	(127, 29), 17	(131, 73), 45	(137, 79), 51	(139, 167), 110	(143, 61), 41															
27	121	(73, 157), 54	(79, 169), 63	(83, 47), 18	(89, 209), 88	(97, 103), 47	(101, 101), 48	(103, 97), 47	(107, 113), 57	(109, 109), 56	(113, 107), 57	(121, 211), 121	(127, 163), 98	(131, 41), 25	(137, 53), 34	(139, 19), 12	(143, 17), 11															
28	127	(73, 19), 6	(79, 193), 72	(83, 209), 82	(89, 143), 60	(97, 181), 83	(101, 47), 22	(103, 199), 97	(107, 11), 5	(109, 163), 84	(113, 29), 15	(121, 67), 38	(127, 211), 127	(131, 17), 10	(137, 191), 124	(139, 13), 8	(143, 89), 60															
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32	143	(73, 71), 24	(79, 47), 17	(83, 151), 59	(89, 37), 15	(97, 179), 82	(101, 43), 20	(103, 191), 93	(107, 19), 9	(109, 167), 86	(113, 31), 16	(121, 173), 99	(127, 59), 35	(131, 163), 101	(137, 139), 90	(139, 137), 90	(143, 211), 143															
33	149	(73, 143), 49	(79, 71), 26	(83, 103), 40	(89, 181), 76	(97, 47), 21	(101, 199), 95	(103, 83), 40	(107, 127), 64	(109, 11), 5	(113, 163), 87	(121, 29), 16	(127, 107), 64	(131, 139), 86	(137, 67), 43	(139, 131), 86	(143, 73), 49															
34	151	(73, 97), 33	(79, 79), 29	(83, 17), 6	(89, 89), 37	(97, 73), 33	(101, 41), 19	(103, 187), 91	(107, 23), 11	(109, 169), 87	(113, 137), 73	(121, 121), 69	(127, 193), 116	(131, 131), 81	(137, 113), 73	(139, 199), 131	(143, 167), 113															
35	157	(73, 169), 58	(79, 103), 38	(83, 179), 70	(89, 23), 9	(97, 151), 69	(101, 197), 94	(103, 79), 38	(107, 131), 66	(109, 13), 6	(113, 59), 31	(121, 187), 107	(127, 31), 18	(131, 107), 66	(137, 41), 26	(139, 193), 127	(143, 29), 19															
36	163	(73, 31), 10	(79, 127), 47	(83, 131), 51	(89, 167), 70	(97, 19), 8	(101, 143), 68	(103, 181), 88	(107, 29), 14	(109, 67), 34	(113, 191), 102	(121, 43), 24	(127, 79), 47	(131, 83), 51	(137, 179), 116	(139, 187), 123	(143, 101), 68															
37	167	(73, 149), 51	(79, 143), 53	(83, 169), 66	(89, 193), 81	(97, 71), 32	(101, 37), 17	(103, 179), 87	(107, 31), 15	(109, 173), 89	(113, 139), 74	(121, 17), 9	(127, 41), 24	(131, 67), 41	(137, 61), 39	(139, 113), 74	(143, 79), 53															
38	169	(73, 103), 35	(79, 151), 56	(83, 83), 32	(89, 101), 42	(97, 97), 44	(101, 89), 42	(103, 73), 35	(107, 137), 69	(109, 121), 62	(113, 113), 60	(121, 109), 62	(127, 127), 76	(131, 59), 36	(137, 107), 69	(139, 181), 119	(143, 173), 117															
39	173	(73, 11), 3	(79, 167), 62	(83, 121), 47	(89, 127), 53	(97, 149), 68	(101, 193), 92	(103, 71), 34	(107, 139), 70	(109, 17), 8	(113, 61), 32	(121, 83), 47	(127, 89), 53	(131, 43), 26	(137, 199), 129	(139, 107), 70	(143, 151), 102															
40	179	(73, 83), 28	(79, 191), 71	(83, 73), 28	(89, 61), 25	(97, 17), 7	(101, 139), 66	(103, 173), 84	(107, 37), 18	(109, 71), 36	(113, 193), 103	(121, 149), 85	(127, 137), 82	(131, 19), 11	(137, 127), 82	(139, 101), 66	(143, 13), 8															
41	181	(73, 37), 12	(79, 199), 74	(83, 197), 77	(89, 179), 75	(97, 43), 19	(101, 191), 91	(103, 67), 32	(107, 143), 72	(109, 19), 9	(113, 167), 89	(121, 31), 17	(127, 13), 7	(131, 11), 6	(137, 173), 112	(139, 169), 111	(143, 107), 72															
42	187	(73, 109), 37	(79, 13), 4	(83, 149), 58	(89, 113), 47	(97, 121), 55	(101, 137), 65	(103, 169), 82	(107, 41), 20	(109, 73), 37	(113, 89), 47	(121, 97), 55	(127, 61), 36	(131, 197), 122	(137, 101), 65	(139, 163), 107	(143, 179), 121															
43	191	(73, 17), 5	(79, 29), 10	(83, 187), 73	(89, 139), 58	(97, 173), 79	(101, 31), 14	(103, 167), 81	(107, 43), 21	(109, 17																						

Continued

		$j$	33	34	35	36	37	38	39	40	41	42	43	44	45	46	47	48
		$q_j$																
$i$	$r_i$		149	151	157	163	167	169	173	179	181	187	191	193	197	199	209	211
1	11	(149, 79), 56	(151, 71), 51	(157, 83), 62	(163, 107), 83	(167, 73), 58	(169, 179), 144	(173, 187), 154	(179, 169), 144	(181, 101), 87	(187, 173), 154	(191, 121), 110	(193, 197), 181	(197, 193), 181	(199, 209), 198	(209, 199), 198	(211, 11), 11	
2	13	(149, 17), 12	(151, 103), 74	(157, 79), 59	(163, 31), 24	(167, 29), 23	(169, 97), 78	(173, 11), 9	(179, 47), 40	(181, 43), 37	(187, 109), 97	(191, 143), 130	(193, 61), 56	(197, 209), 196	(199, 37), 35	(209, 197), 196	(211, 13), 13	
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4	19	(149, 41), 29	(151, 199), 143	(157, 67), 50	(163, 13), 10	(167, 107), 85	(169, 61), 49	(173, 113), 93	(179, 101), 86	(181, 79), 68	(187, 127), 113	(191, 209), 190	(193, 73), 67	(197, 47), 44	(199, 151), 143	(209, 191), 190	(211, 19), 19	
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8	37	(149, 113), 80	(151, 67), 48	(157, 31), 23	(163, 169), 131	(167, 131), 104	(169, 163), 131	(173, 209), 172	(179, 53), 45	(181, 187), 161	(187, 181), 161	(191, 197), 179	(193, 109), 100	(197, 191), 179	(199, 73), 69	(209, 173), 172	(211, 37), 37	
9	41	(149, 199), 141	(151, 131), 94	(157, 23), 17	(163, 17), 13	(167, 43), 34	(169, 209), 168	(173, 67), 55	(179, 19), 16	(181, 71), 61	(187, 53), 47	(191, 31), 28	(193, 47), 43	(197, 13), 12	(199, 149), 141	(209, 169), 168	(211, 41), 41	
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11	47	(149, 13), 9	(151, 17), 12	(157, 11), 8	(163, 209), 162	(167, 121), 96	(169, 173), 139	(173, 169), 139	(179, 73), 62	(181, 107), 92	(187, 71), 63	(191, 97), 88	(193, 59), 54	(197, 61), 57	(199, 53), 50	(209, 163), 162	(211, 47), 47	
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20	89	(149, 181), 128	(151, 59), 42	(157, 137), 102	(163, 83), 64	(167, 37), 29	(169, 131), 105	(173, 43), 35	(179, 31), 26	(181, 149), 128	(187, 197), 175	(191, 139), 126	(193, 143), 131	(197, 187), 175	(199, 11), 10	(209, 121), 120	(211, 89), 89	
21	97	(149, 143), 101	(151, 187), 134	(157, 121), 90	(163, 199), 154	(167, 71), 56	(169, 13), 10	(173, 179), 147	(179, 173), 147	(181, 127), 109	(187, 151), 134	(191, 17), 15	(193, 19), 17	(197, 41), 38	(199, 163), 154	(209, 113), 112	(211, 97), 97	
22	101	(149, 19), 13	(151, 41), 29	(157, 113), 84	(163, 47), 36	(167, 193), 153	(169, 59), 47	(173, 37), 30	(179, 139), 118	(181, 11), 9	(187, 23), 20	(191, 61), 55	(193, 167), 153	(197, 73), 68	(199, 29), 27	(209, 109), 108	(211, 101), 101	
23	103	(149, 167), 118	(151, 73), 52	(157, 109), 81	(163, 181), 140	(167, 149), 118	(169, 187), 150	(173, 71), 58	(179, 17), 14	(181, 163), 140	(187, 169), 150	(191, 83), 75	(193, 31), 28	(197, 89), 83	(199, 67), 63	(209, 107), 106	(211, 103), 103	
24	107	(149, 43), 30	(151, 137), 98	(157, 101), 75	(163, 29), 22	(167, 61), 48	(169, 23), 18	(173, 139), 114	(179, 193), 164	(181, 47), 40	(187, 41), 36	(191, 127), 115	(193, 179), 164	(197, 121), 113	(199, 143), 135	(209, 103), 102	(211, 107), 107	
25	109	(149, 191), 135	(151, 169), 121	(157, 97), 72	(163, 163), 126	(167, 17), 13	(169, 151), 121	(173, 173), 142	(179, 71), 60	(181, 199), 171	(187, 187), 166	(191, 149), 135	(193, 43), 39	(197, 137), 128	(199, 181), 171	(209, 101), 100	(211, 109), 109	
26	113	(149, 67), 47	(151, 23), 16	(157, 89), 66	(163, 11), 8	(167, 139), 110	(169, 197), 158	(173, 31), 25	(179, 37), 31	(181, 83), 71	(187, 59), 52	(191, 193), 175	(193, 191), 175	(197, 169), 158	(199, 47), 44	(209, 97), 96	(211, 113), 113	
27	121	(149, 29), 20	(151, 151), 108	(157, 73), 54	(163, 127), 98	(167, 173), 137	(169, 79), 63	(173, 167), 137	(179, 179), 152	(181, 61), 52	(187, 13), 11	(191, 71), 64	(193, 67), 61	(197, 23), 21	(199, 199), 188	(209, 89), 88	(211, 121), 121	
28	127	(149, 53), 37	(151, 37), 26	(157, 61), 45	(163, 109), 84	(167, 41), 32	(169, 43), 34	(173, 59), 48	(179, 23), 19	(181, 97), 83	(187, 31), 27	(191, 137), 124	(193, 79), 72	(197, 71), 66	(199, 103), 97	(209, 83), 82	(211, 127), 127	
29	131	(149, 139), 98	(151, 101), 72	(157, 53), 39	(163, 167), 129	(167, 163), 129	(169, 89), 71	(173, 127), 104	(179, 199), 169	(181, 191), 164	(187, 113), 100	(191, 181), 164	(193, 17), 15	(197, 103), 96	(199, 179), 169	(209, 79), 78	(211, 131), 131	
30	137	(149, 163), 115	(151, 197), 141	(157, 41), 30	(163, 149), 115	(167, 31), 24	(169, 53), 42	(173, 19), 15	(179, 43), 36	(181, 17), 14	(187, 131), 116	(191, 37), 33	(193, 29), 26	(197, 151), 141	(199, 83), 78	(209, 73), 72	(211, 137), 137	
31	139	(149, 101), 71	(151, 19), 13	(157, 37), 27	(163, 73), 56	(167, 197), 156	(169, 181), 145	(173, 53), 43	(179, 131), 111	(181, 169), 145	(187, 67), 59	(191, 59), 53	(193, 103), 94	(197, 167), 156	(199, 121), 114	(209, 71), 70	(211, 139), 139	
32	143	(149, 187), 132	(151, 83), 59	(157, 29), 21	(163, 131), 101	(167, 109), 86	(169, 17), 13	(173, 121), 99	(179, 97), 82	(181, 53), 45	(187, 149), 132	(191, 103), 93	(193, 41), 37	(197, 199), 186	(199, 197), 186	(209, 67), 66	(211, 143), 143	
33	149	(149, 211), 149	(151, 179), 128	(157, 17), 12	(163, 113), 87	(167, 187), 148	(169, 191), 153	(173, 13), 10	(179, 151), 128	(181, 89), 76	(187, 167), 148	(191, 169), 153	(193, 53), 48	(197, 37), 34	(199, 101), 95	(209, 61), 60	(211, 149), 149	
34	151	(149, 149), 105	(151, 211), 151	(157, 13), 9	(163, 37), 28	(167, 143), 113	(169, 109), 87	(173, 47), 38	(179, 29), 24	(181, 31), 26	(187, 103), 91	(191, 191), 173	(193, 127), 116	(197, 53), 49	(199, 139), 131	(209, 59), 58	(211, 151), 151	
35	157	(149, 173), 122	(151, 97), 69	(157, 211), 157	(163, 19), 14	(167, 11), 8	(169, 73), 58	(173, 149), 122	(179, 83), 70	(181, 67), 57	(187, 121), 107	(191, 47), 42	(193, 139), 127	(197, 101), 94	(199, 43), 40	(209, 53), 52	(211, 157), 157	
36	163	(149, 197), 139	(151, 193), 138	(157, 199), 148	(163, 211), 163	(167, 89), 70	(169, 37), 29	(173, 41), 33	(179, 137), 116	(181, 103), 88	(187, 139), 123	(191, 113), 102	(193, 151), 138	(197, 149), 139	(199, 157), 148	(209, 47), 46	(211, 163), 163	
37	167	(149, 73), 51	(151, 47), 33	(157, 191), 142	(163, 59), 45	(167, 211), 167	(169, 83), 66	(173, 109), 89	(179, 103), 87	(181, 197), 169	(187, 11), 9	(191, 157), 142	(193, 89), 81	(197, 181), 169	(199, 23), 21	(209, 43), 42	(211, 167), 167	
38	169	(149, 11), 7	(151, 79), 56	(157, 187), 139	(163, 193), 149	(167, 167), 132	(169, 211), 169	(173, 143), 117	(179, 191), 162	(181, 139), 119	(187, 157), 139	(191, 179), 162	(193, 163), 149	(197, 197), 184	(199, 61), 57	(209, 41), 40	(211, 169), 169	
39	173	(149, 97), 68	(151, 143), 102	(157, 179), 133	(163, 41), 31	(167, 79), 62	(169, 47), 37	(173, 211), 173	(179, 157), 133	(181, 23), 19	(187, 29), 25	(191, 13), 11	(193, 101), 92	(197, 19), 17	(199, 137), 129	(209, 37), 36	(211, 173), 173	
40	179	(149, 121), 85	(151, 29), 20	(157, 167), 124	(163, 23), 17	(167, 157), 124	(169, 11), 8	(173, 103), 84	(179, 211), 179	(181, 59), 50	(187, 47), 41	(191, 79), 71	(193, 113), 103	(197, 67), 62	(199, 41), 38	(209, 31), 30	(211, 179), 179	
41	181	(149, 59), 41	(151, 61), 43	(157, 163), 121	(163, 157), 121	(167, 113), 89	(169, 139), 111	(173, 137), 112	(179, 89), 75	(181, 211), 181	(18							

3) If  $(q_j, q_{j\hat{i}})$  is not a co-factor-pair, then  $(q_{j\hat{i}}, q_j)$  is a factor-pair with the same generation  $t_{j\hat{i}}$  as  $(q_j, q_{j\hat{i}})$ .

4) For any  $j$  and  $i$ , we have  $t_{j\hat{i}} \leq r_j$ .

We found that the factors of a composite number  $\alpha$  are closely related to factor-pairs in the Factor-Pair Table with generation information:

**Theorem 2.** If  $q_j + 210\theta$  ( $\theta \in N_0$ ) is a factor of  $\alpha = r_i + 210k$ , then  $\frac{\alpha}{q_j + 210\theta}$  is a descendant of  $q_{j\hat{i}}$ , i.e.,

$$\frac{\alpha}{q_j + 210\theta} = q_{j\hat{i}} + 210\hat{\theta}. \tag{3}$$

Theorems 1 and 2 confirm that any composite number  $\alpha$  without factors of 2, 3, 5 and 7 can be written as

$$\alpha = r_i + 210k = (q_j + 210\theta)(q_{j\hat{i}} + 210\hat{\theta}), \tag{4}$$

where  $\theta, \hat{\theta} \in N_0$  and  $j \in \{1, 2, \dots, 48\}$ .

From Equation (4), we have

$$r_i + 210k = q_j q_{j\hat{i}} + 210\theta q_{j\hat{i}} + 210\hat{\theta} q_j + 210^2 \theta \hat{\theta}. \tag{5}$$

Combining Equations (2) and (5), we get

$$t_{j\hat{i}} = k - \theta q_{j\hat{i}} - \hat{\theta} q_j - 210\theta \hat{\theta}. \tag{6}$$

When  $j$  is fixed, Equation (6) is a quadratic equation with only two unknowns  $\theta$  and  $\hat{\theta}$  in  $N_0$ .

**Theorem 3.** To check whether a number  $\alpha = r_i + 210k$  is a composite number, it suffices to check if Equation (6) has a solution  $(\theta, \hat{\theta}) \in N_0 \times N_0$  for some  $j \in \{1, 2, \dots, 48\}$ .

We then propose an algorithm to solve the equation.

Without loss of generality, we assume  $q_j + 210\theta \leq q_{j\hat{i}} + 210\hat{\theta}$ , since we can replace  $q_j$  and  $\theta$  with  $q_{j\hat{i}}$  and  $\hat{\theta}$  correspondingly if  $q_j + 210\theta > q_{j\hat{i}} + 210\hat{\theta}$ . Given  $\alpha = (q_j + 210\theta)(q_{j\hat{i}} + 210\hat{\theta})$ , we see

$$q_j + 210\theta \leq \sqrt{\alpha}. \tag{7}$$

Therefore,

$$\theta \leq \frac{\sqrt{\alpha} - q_j}{210} \leq \frac{\sqrt{\alpha}}{210}. \tag{8}$$

We can derive a direct corollary from Theorem 3:

**Corollary 1.** Checking whether a number  $\alpha = r_i + 210k \in \Lambda$  is a composite number is equivalent to checking if

$$\hat{\theta} = \frac{k - t_{j\hat{i}} - \theta q_{j\hat{i}}}{q_j + 210\theta} \tag{9}$$

is an integer for some  $\theta \in \left\{0, 1, \dots, \left\lceil \frac{\sqrt{\alpha}}{210} \right\rceil\right\}$ .

Based on this Corollary, we can develop a new algorithm for primality test:

**Algorithm 1:**

Step 1: Input a natural number  $\alpha$ . If  $\alpha$  is divisible by 2, 3, 5 or 7, then  $\alpha$  is not a prime, and the algorithm terminates.

Step 2: Set  $r_i = \alpha - \left\lfloor \frac{\alpha - 10}{210} \right\rfloor \times 210$ .

Step 3: For each  $j \in \{1, 2, \dots, 48\}$ , examine the Factor-Pair Table with generation information for the factor pair  $(q_j, q_{j_i})$  and generation  $t_{j_i}$ . For each

$\theta \in \left\{ 0, 1, \dots, \left\lfloor \frac{\sqrt{\alpha}}{210} \right\rfloor \right\}$ , compute Equation (9). According to Corollary 1, if  $\hat{\theta} \in N_0$ , then

$\alpha$  is not a prime, and the algorithm terminates.

Step 4: If the loop completes (having traversed all  $j$ ) without finding a solution, then  $\alpha$  is a prime.

The following theorem gives the complexity of Algorithm 1:

**Theorem 4.** Algorithm 1 involves at most 48 linear searches over  $\left\lfloor \frac{\sqrt{\alpha}}{210} \right\rfloor + 1$

integer values. In other words, the worst case is searching through about  $\sqrt{\alpha}/4$  integer values.

However, to list all primes within a pre-specified interval, the above algorithm for primality test costs too much time. Therefore, we need to find a more efficient way to locate primes.

### 3. Method

Notice that in a given interval, primes only appear with the form of  $r_i + 210k$ , where  $r_i \in S$  and  $k \in N_0$ , since the rest is divisible by 2, 3, 5 or 7. So we only need to consider the primality of those numbers.

Furthermore, in our PTP method to list primes in the following observations, the cyclic behavior of the Factor-Pair Table with generation information is observed and utilized to ease computation burden.

#### 3.1. Basic Binary Matrices

According to the 4th property of the Factor-Pair Table with generation information, we know

$$t_{j_i} + r_j \leq 2r_j. \tag{10}$$

To represent the generation information of each factor pair more clearly, we convert each column of the Factor-Pair Table with generation information,  $V_j$ , into a basic binary matrix of composites  $B_j^0 = [b_j^0(i, k)]_{48 \times 2r_j}$  for  $i = 1, 2, \dots, 48$ ,  $k = 0, 1, 2, \dots, 2r_j - 1$ , where  $b_j^0(i, k) \in \{0, 1\}$ . In this matrix,  $b_j^0(i, k) = 1$  indicates that the  $k^{\text{th}}$  descendant of  $r_i$  is a composite with a factor being  $q_j$ .

Specifically, we specify  $b_j^0(i, k)$  following the rule below:

The entry  $b_j^0(i, k)$  is set as 1 only under two cases:

Case 1 ( $i = j$ ): if  $k = r_i$  or  $k = 0$ ,  $b_j^0(i, k) = 1$ .

Case 2 ( $i \neq j$ ): if  $k = t_{ji}$  or  $k = t_{ji} + r_j$ ,  $b_j^0(i, k) = 1$ .

Otherwise,  $b_j^0(i, k) = 0$ .

Inequality (10) guarantees that the rule is properly defined.

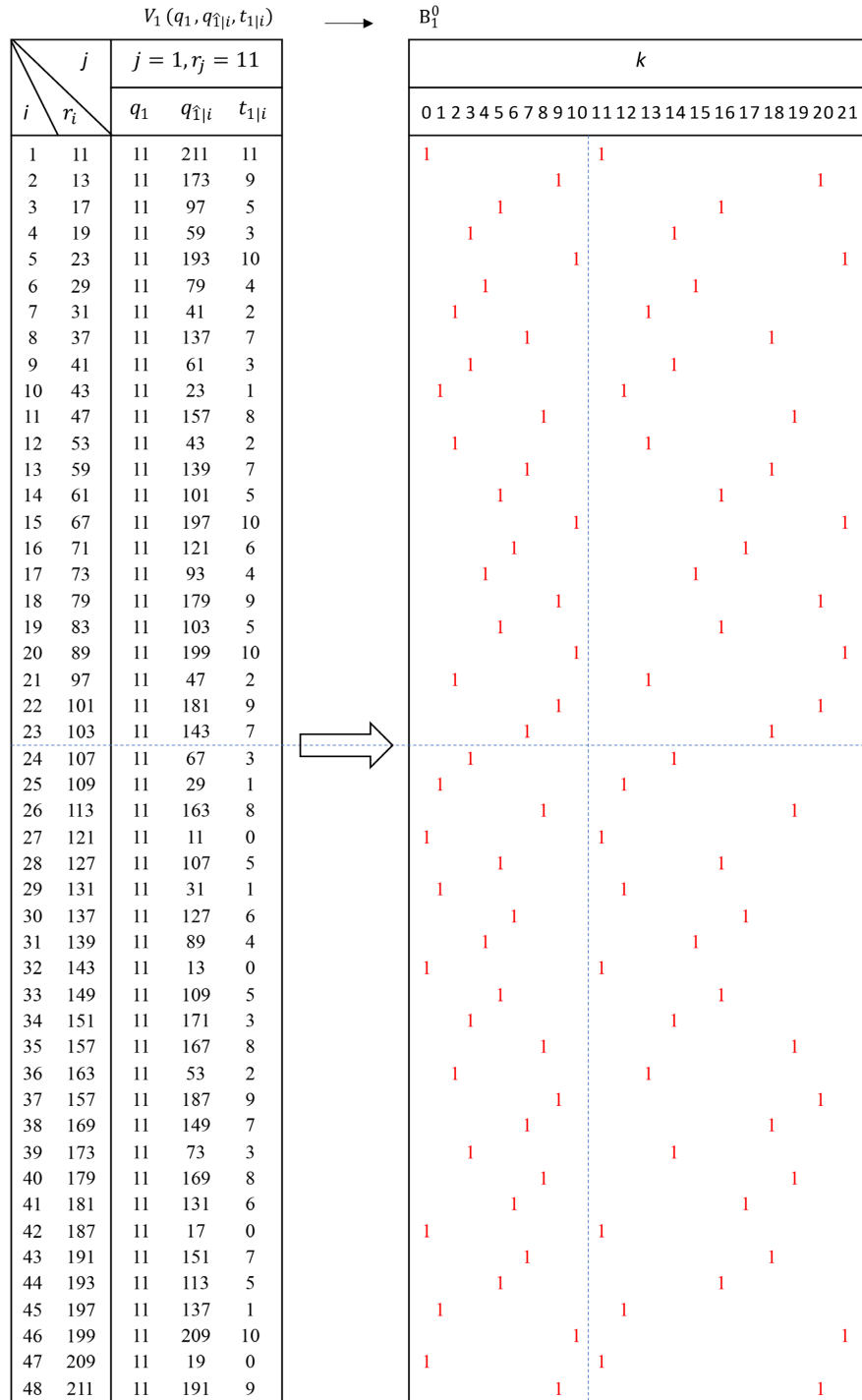


Figure 1. Converting  $V_1$  to  $B_1^0$ .

We observe the following properties in  $B_j^0$ :

1) There are exactly two non-zero entries for each  $i$ , which means either Case 1 or Case 2 always sets two entries in  $B_j^0$  as 1.

2) Every matrix  $B_j^0$  consists of two identical submatrices of the size  $48 \times r_j$  placed side by side.

3) Each non-zero entry in the submatrix mentioned in (2) corresponds to a factor-pair  $(q_j, q_{j\bar{i}})$ ,  $i = 1, 2, \dots, 48$  with its generation  $t_{j\bar{i}}$ .

**Figure 1** shows how to convert  $V_1$  to  $B_1^0$ . For instance,

we set  $b_1^0(1, 0) = b_1^0(1, 11) = 1$ , moreover,

since  $t_{1|2} = 9$ ,  $b_1^0(2, 9) = b_1^0(2, 20) = 1$ ;

since  $t_{1|3} = 5$ ,  $b_1^0(3, 5) = b_1^0(3, 16) = 1$ ;

since  $t_{1|4} = 3$ ,  $b_1^0(4, 3) = b_1^0(4, 14) = 1$ .

Similarly,  $V_2, V_3, \dots$  and  $V_{48}$  is respectively converted into  $B_2^0, B_3^0, \dots$  and  $B_{48}^0$ .

### 3.2. Mapping Effect and Refined Binary Matrices

From Sections 2.1 and 2.2, we know a composite without factors of 2, 3, 5 and 7, which is rooted as  $r_i$ , can be expressed as

$\alpha = r_i + 210k = (q_j + 210\theta) \times (q_{j\bar{i}} + 210\hat{\theta})$ , and its generation is

$$t_{j\bar{i}} = \frac{q_j q_{j\bar{i}} - r_i}{210} \in N_0.$$

Now if  $\theta = 0$ , we observe that for  $k = t_{j\bar{i}} + q_j \times \hat{\theta}$ , where  $\hat{\theta} \in N_0$ ,  $r_i + 210t_{j\bar{i}}$ ,  $r_i + 210(t_{j\bar{i}} + q_j)$ ,  $r_i + 210(t_{j\bar{i}} + 2q_j)$ ,  $r_i + 210(t_{j\bar{i}} + 3q_j), \dots$ , are all composites rooted at  $r_i$ . That means, there is a mapping effect between  $r_i + 210t_{j\bar{i}}$  and  $r_i + 210(t_{j\bar{i}} + m \times q_j)$  for any  $m \in N_+$ .

Similarly, if  $\hat{\theta} = 0$ , there is a mapping effect between  $r_i + 210t_{j\bar{i}}$  and  $r_i + 210(t_{j\bar{i}} + m' \times q_{j\bar{i}})$  for  $m' \in N_+$ . Such an effect is very useful to quickly find the same-root composites based on the basic-binary-matrices, as described below:

Consider an integer interval from  $11 + 210 \times k^*$  to  $211 + 210 \times (k^* + \delta)$ , where  $k^* \in N_+$ ,  $\delta \in N_0$ ,  $\delta \leq 10 (< r_1 = 11)$ . Define  $B_j = [b_j(i, k^* + h)]_{48 \times (\delta + 1)}$ , where  $b_j(i, k^* + h) \in \{0, 1\}$ , as a refined binary matrix of composites for  $i \in \{1, 2, \dots, 48\}$ ,  $h \in \{0, 1, \dots, \delta\}$ . In other words, if  $b_j(i, k) = 1$ , then the integer  $r_i + 210k$  is a composite number. But if  $b_j(i, k) = 0$ , we can't determine whether  $r_i + 210k$  is a prime, as discussed later. Theorem 5 then shows the mapping effect from  $B_j^0$  to  $B_j$ :

**Theorem 5.** If  $k = w_j + r_j \times M$ , where  $M \in N_0$  and  $w_j \in \{0, 1, \dots, 2r_j - 1\}$ , then for any  $i \in \{1, 2, \dots, 48\}$ , if  $b_j^0(i, w_j) = 1$ , then  $r_i + 210k$  is composite, i.e.,  $b_j(i, k) = 1$ .

The value of each entry is determined accordingly. By setting  $w_j = k^* - r_j \lfloor k^*/r_j \rfloor$ , we notice that  $w_j + h \in \{0, 1, \dots, 2r_j - 1\}$ . For any  $h \in \{0, 1, \dots, \delta\}$  and  $i \in \{1, 2, \dots, 48\}$ , if  $b_j^0(i, w_j + h) = 1$ , then set  $b_j(i, k^* + h) = 1$ ; if  $b_j^0(i, w_j + h) = 0$ , then set  $b_j(i, k^* + h) = 0$ .

$B_j$  has the following properties:

- 1)  $B_j$  is a submatrix of  $B_j^0$ , with dimensionality of  $48 \times (\delta + 1)$ .
- 2) Each non-zero entry in  $B_j$  corresponds to a composite number in the pre-specified interval.

The refined binary matrix  $B_j$  can quickly locate large composites in a pre-specified interval. However, it's not guaranteed that when  $b_j(i, k^* + h) = 0$ , the corresponding integer is prime. The reason is that the mapping effect only locates composites with the form of  $r_i + 210(t_{j_i} + m \times q_j) = q_j \times (q_{j_i} + 210m)$  and  $r_i + 210(t_{j_i} + m' \times q_{j_i}) = q_{j_i} \times (q_j + 210m')$ . For composites with the smallest factor larger than 211,  $B_j$  may fail to detect them.

### 3.3. Composite Binary Matrix

A composite binary matrix  $C = [c(i, k)]_{48 \times (\delta + 1)}$ , where  $i \in \{1, 2, \dots, 48\}$  and  $k \in \{k^*, k^* + 1, \dots, k^* + \delta\}$ , can be formed based on  $B_1, B_2, \dots, B_{48}$ . We want to assure that all composites within the interval  $[11 + 210 \times k^*, 211 + 210 \times (k^* + \delta)]$  are detected and located by  $C$ .

We first locate composite numbers detected by mapping effect. If  $b_j(i, k) = 1$  for some  $j \in \{1, 2, \dots, 48\}$ , the number  $r_i + 210k$  is a composite, i.e.,  $c(i, k) = 1$ . After that, we only need to determine the primality of the remaining numbers.

The specific algorithm is as follows:

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#### Algorithm 2:

Step 1: For each  $c(i, k)$ , where  $i \in \{1, 2, \dots, 48\}$  and  $k \in \{k^*, k^* + 1, \dots, k^* + \delta\}$ , check  $b_j(i, k)$  for  $j \in \{1, 2, \dots, 48\}$ :

If  $b_j(i, k) = 1$  for some  $j$ , then set  $c(i, k) = 1$ ;  
 otherwise, go to Step 2.

Step 2: Let  $\alpha = r_i + 210k$ . Conduct a selected primality test algorithm for  $\alpha$ :

If  $\alpha$  is not a prime, set  $c(i, k) = 1$ ;  
 otherwise, set  $c(i, k) = 0$ .

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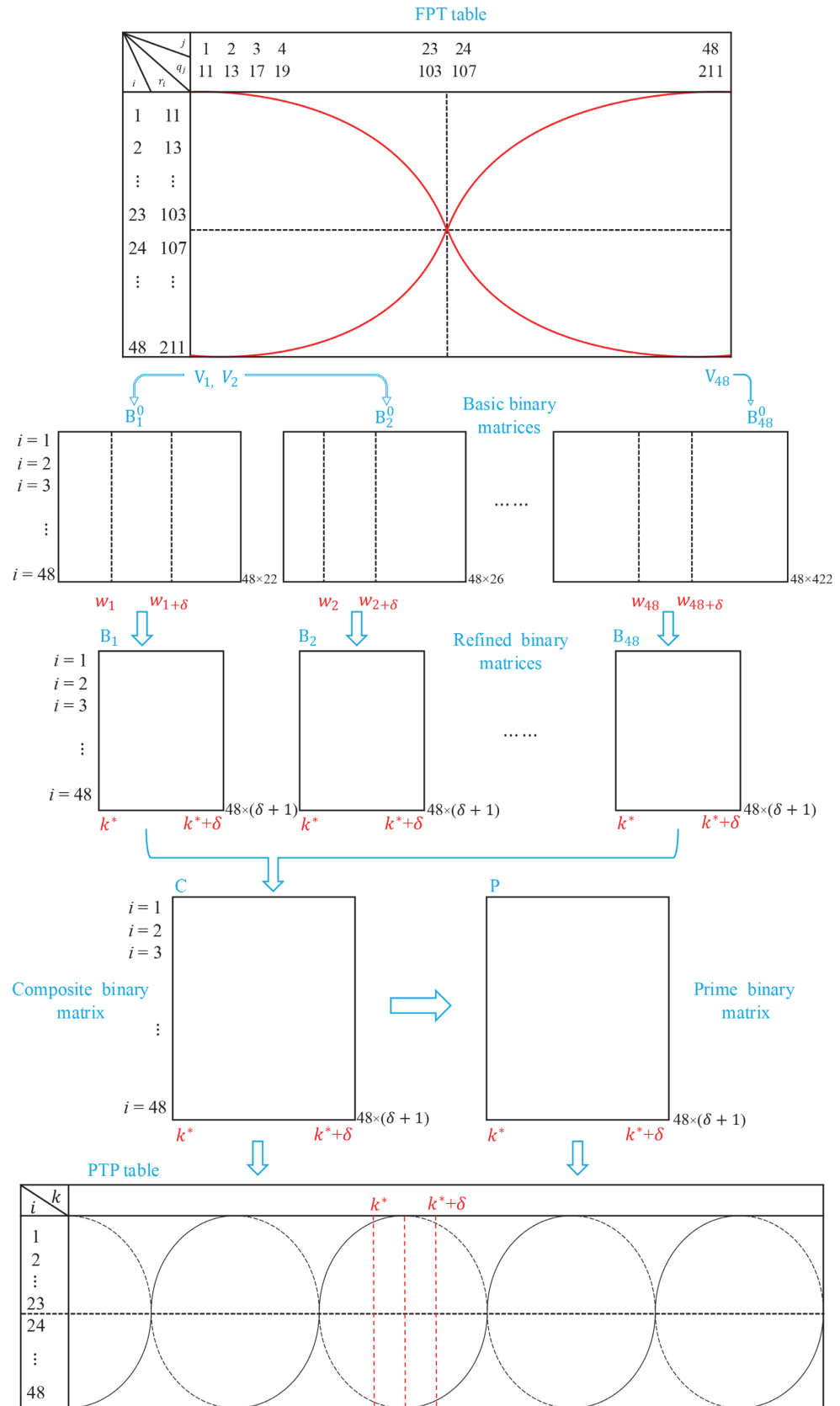
With matrix  $C$ , we can locate all composite numbers within the interval:

**Theorem 6.** In a prespecified interval  $[11 + 210 \times k^*, 211 + 210 \times (k^* + \delta)]$ , where  $k^* \in N_+$ ,  $\delta \in [0, 10]_N$ , any integer number  $r_i + 210k$  is composite if and only if  $c(i, k) = 1$ .

### 3.4. Prime Binary Matrix

We are now able to form a binary matrix of primes  $P = [p(i, k)]_{48 \times (\delta + 1)}$ , where  $i \in \{1, 2, \dots, 48\}$  and  $k \in \{k^*, k^* + 1, \dots, k^* + \delta\}$ , within the pre-specified interval. Each entry of  $P$  is determined by  $p(i, k) = 1 - c(i, k)$ ,  $i = 1, 2, \dots, 48$  and  $k = k^*, k^* + 1, \dots, k^* + \delta$ .

In this way, we can list all primes within the interval as illustrated in the following theorem:



**Figure 2.** An outline of observations.

**Theorem 7.** In a pre-specified interval  $[11+210 \times k^*, 211+210 \times (k^* + \delta)]$ , where  $k^* \in N_+$ ,  $\delta \in [0, 10]_N$ , any integer number  $r_i + 210k$  is prime if and only if  $p(i, k) = 1$ .

**Figure 2** is an outline of our PTP method for listing all primes within the interval  $[11+210 \times k^*, 211+210 \times (k^* + \delta)]$  for  $\delta \in [0, 10]_N$ .

We first form a  $48 \times 48$  Factor-Pair Table with generation information, then convert each column of it into a basic binary matrix  $B_j^0$ . After that,  $B_j^0$  is modified as a refined binary matrix  $B_j$ . By aggregating  $B_1, B_2, \dots, B_{48}$ , and testing the primality of the rest numbers, we obtain a composite binary matrix  $C$ . Based on  $C = [c(i, k)]_{48 \times (\delta+1)}$ , we have a prime binary matrix  $P = [p(i, k)]_{48 \times (\delta+1)}$ . The primes within the interval are listed as  $\alpha = r_i + 210k$  for  $k \in \{k^*, k^* + 1, \dots, k^* + \delta\}$  such that  $p(i, k) = 1$  in the prime binary matrix  $P$ .

Moreover, all composites (without factors of 2, 3, 5 and 7) within the same interval can be predicted. Both primes and composites are illustrated on a PTP table.

## 4. Results

### 4.1. Formula of Primes

From Theorem 7, we have the following Formula of Primes:

**Formula of Primes.** All prime numbers within an integer interval  $[11+210 \times k^*, 211+210 \times (k^* + \delta)]$ , where  $k^* \in N_+$ ,  $\delta \in [0, 10]_N$ , have the form of  $\alpha = r_i + 210k$ , where  $k \in \{k^*, k^* + 1, \dots, k^* + \delta\}$  and  $i \in \{1, 2, \dots, 48\}$ , such that  $p(i, k) = 1$  in the prime binary matrix  $P$ .

### 4.2. Algorithm of Prime Listing

Finally, an algorithm for listing all primes in a pre-specified interval is presented below.

#### Algorithm 3:

Part A: Data Preparation

Firstly, we need to construct a  $48 \times 48$  Factor-Pair Table with generation information,  $(V_1, V_2, \dots, V_{48})$ . Then we can convert each  $V_j$  into a basic binary matrix  $B_j^0 = (b_j^0(i, k))_{48 \times 2r_j}$ , for  $i = 1, 2, \dots, 48$  and  $k = 0, 1, 2, \dots, 2r_j - 1$ .

Part B: Prime Prediction (Given  $k^* \in N_+$ ,  $\delta \in [0, 10]_N$ )

Step 1: Generate refined binary matrices  $B_j = (b_j(i, k))_{48 \times (\delta+1)}$  for  $j = 1, 2, \dots, 48$

based on  $B_j^0$ .

Step 2: Generate a composite binary matrix  $C = (c(i, k))_{48 \times (\delta+1)}$  with Algorithm 2.

Step 3: Generate a prime binary matrix  $P = (p(i, k))_{48 \times (\delta+1)}$  based on  $C$ .

Step 4: The prime numbers within the interval  $[11+210 \times k^*, 211+210 \times (k^* + \delta)]$  can be listed as

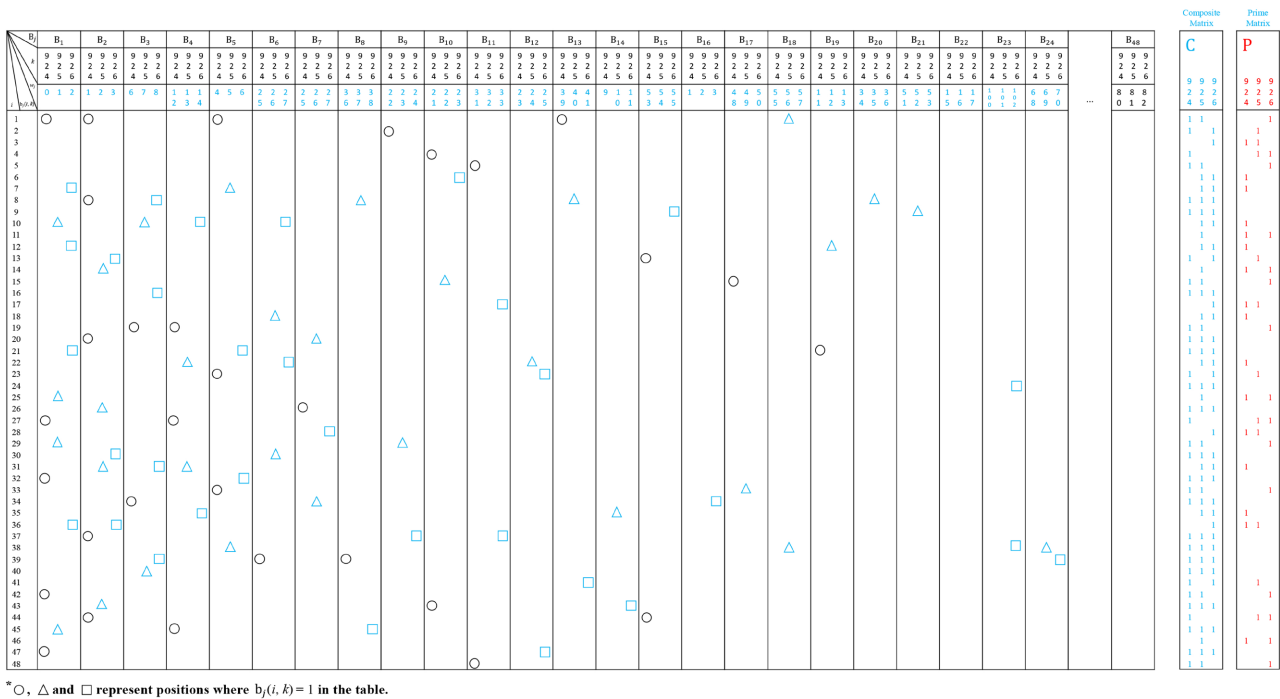
$$\{\alpha = r_i + 210k \mid p(i, k) = 1, i = 1, 2, \dots, 48, k = k^*, k^* + 1, \dots, k^* + \delta\}.$$

In the above algorithm, Part A and Step 1, Step 3 and Step 4 in Part B require constant time complexity. Step 2 in Part B involves a few primality tests, whose complexity can be reduced by choosing a more efficient primality test algorithm.

An example of Algorithm 3 can be found in Appendix A.

The procedure of Part B is presented in **Figure 3**. The top row shows  $B_1, B_2, \dots, B_{48}, C$  and  $P$ , respectively. The second row is for  $k^* = 924$ ,  $k^* + 1 = 925$  and  $k^* + 2 = 926$ . The third row is for  $(\delta_j, \delta_j + 1, \delta_j + 2)$ , where  $j = 1, 2, \dots, 48$ . The  $b_j(i, k)$  entry denotes the binary value for given  $B_j, r_i$  and  $k$ .

A table of listing all primes in  $[11 + 210 \times 922, 211 + 210 \times 928]$  is shown in **Table 2**.



**Figure 3.** Process of predicting primes in  $[11 + 210 \times 924, 211 + 210 \times 926]$ .

**Table 2.** Prediction of Primes within  $[11 + 210 \times 922, 211 + 210 \times 928]$ .

$i$	$k$	924	925	926
1	11			194,471
2	13		194,263	
3	17	194,057	194,267	
4	19		194,269	194,479
5	23			194,483
6	29	194,069		

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**Continued**

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7	31	194,071		
8	37			
9	41			
10	43	194,083		
11	47	194,087		194,507
12	53	194,093		
13	59		194,309	
14	61	194,101		194,521
15	67			194,527
16	71			
17	73	194,113	194,323	
18	79	194,119		
19	83			194,543
20	89			
21	97			
22	101	194,141		
23	103		194,353	
24	107			
25	109	194,149		194,569
26	113			
27	121		194,371	194,581
28	127	194,167	194,377	
29	131			194,591
30	137			
31	139	194,179		
32	143			
33	149			194,609
34	151			
35	157	194,197		
36	163	194,203	194,413	
37	167			
38	169			

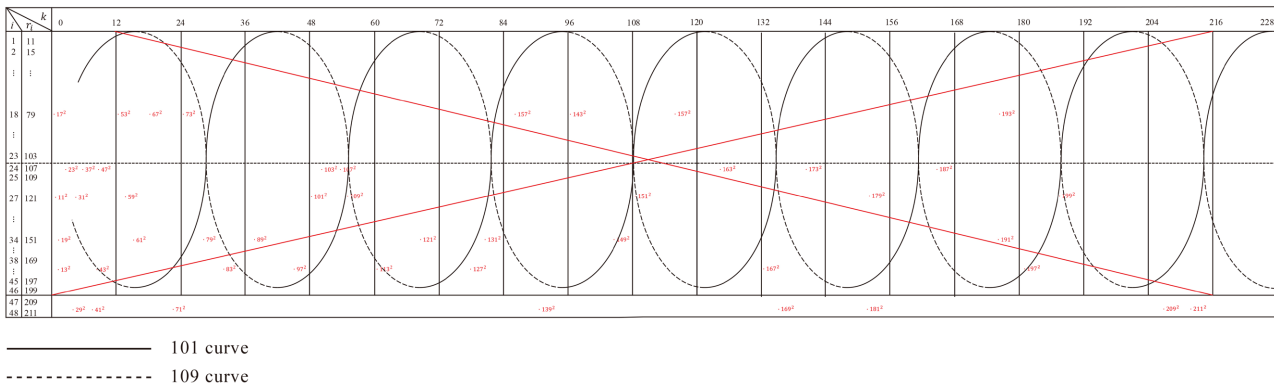
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Continued

39	173		
40	179		
41	181	194,431	
42	187		194,647
43	191		
44	193	194,443	194,653
45	197		
46	199	194,239	194,659
47	209		
48	211		194,671

### 4.3. The Periodic Table of Primes

The Periodic Table of Primes (PTP) is shown in **Figure 4**. Reconstructed from our prime listing algorithm, it illustrates that composites (without factors of 2, 3, 5 and 7) are distributed cyclically and primes are distributed periodically.



**Figure 4.** PTP [11, 211 + 210<sup>2</sup>].

### 5. Discussion

In this paper, we propose a novel way to list primes. Given an interval, without loss of generality, we first restrict the range of numbers to those having the form of  $\alpha = r_i + 210k$ . Then, by mapping the effect, we further eliminate a good amount of numbers and conduct primality checking on the rest of the numbers. Subsequently, we list all primes in the given interval.

Some further discussions are as follows:

1) With theoretical analysis and computer validation (in Appendix B), we can, on average, identify about 26 composite numbers in a column. This demonstrates that the proposed method largely reduces the number of primality checks required in the prime listing task over a pre-specified interval with simple mappings. The

details are found in Appendix B.

2) In practice, we can use bitwise operations to speed up the algorithm. We first encode each column of the basic binary matrices into a binary number, and the same applies to the columns of the refined binary matrices. Then, instead of checking individual entries, we perform the bitwise operation “OR” on all columns in the refined binary matrix corresponding to a certain column. With a significant improvement in efficiency, the resulting value implies the composites located in this column by Step 1 of Algorithm 2.

3) In this paper and all associated demonstrations, we consider composites without factors of 2, 3, 5 and 7 to generate a Factor-Pair Table which gives way to a designated and graphics-manageable period of  $2 \times 3 \times 5 \times 7 = 210$ . To improve efficiency and to list a larger number of primes, we may consider composites without factors of 2, 3, 5, 7 and 11 which represent a larger period of  $2 \times 3 \times 5 \times 7 \times 11 = 2310$ . Further extension is also possible for even larger dimension by adding primes 13, 17, ...

4) The proposed algorithm is obviously capable of listing prime numbers within a specific interval starting from a very large number. Unlike the conventional methods, the proposed algorithm does not require prior knowledge of smaller prime numbers. Furthermore, the mapping process has constant time and small space complexity, ensuring high efficiency.

5) We are investigating the application of the proposed method in various domains, such as cryptography-security, hash-data-base-design and DNA-sequence-identify.

## Conflicts of Interest

The authors declare no conflicts of interest regarding the publication of this paper.

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### Appendix A: An Example of Algorithm 3

**Example:** (List all primes within the interval  $[11+210 \times 924, 211+210 \times 926]$ )

Part A:

The Factor-Pair Table with generation information is shown in **Table 1** and the basic binary matrices  $B_j^0$ ,  $j=1, 2, \dots, 48$ , are converted accordingly.

Part B:

Step 1: Generate refined binary matrices  $B_j = (b_j(i, k))_{48 \times 3}$  for  $j=1, 2, \dots, 48$  based on  $B_j^0$ .

Take  $j=1$  for example,  $B_1 = (b_1(i, k))_{48 \times 3}$ , where  $i=1, 2, \dots, 48$  and  $k=924, 925, 926$ , and  $w_1 = 924 - 11 \times \lfloor 924/11 \rfloor = 0$ . From Theorem 5,

$b_1(i, 924+h) = b_1^0(i, w_1+h) = b_1^0(i, h)$ . From  $B_1^0$  in **Figure 1**, we know that

$$\therefore b_1^0(1, 0) = b_1^0(27, 0) = b_1^0(32, 0) = b_1^0(42, 0) = b_1^0(47, 0) = 1,$$

$$\therefore b_1(1, 924) = b_1(27, 924) = b_1(32, 924) = b_1(42, 924) = b_1(47, 924) = 1.$$

Similarly,

$$\therefore b_1^0(10, 1) = b_1^0(25, 1) = b_1^0(29, 1) = b_1^0(45, 1) = 1,$$

$$\therefore b_1(10, 925) = b_1(25, 925) = b_1(29, 925) = b_1(45, 925) = 1.$$

$$\therefore b_1^0(7, 2) = b_1^0(12, 2) = b_1^0(21, 2) = b_1^0(36, 2) = 1,$$

$$\therefore b_1(7, 926) = b_1(12, 926) = b_1(21, 926) = b_1(36, 926) = 1.$$

Step 2: Generate a composite binary matrix  $C = (c(i, k))_{48 \times 3}$  with Algorithm 2.

Check  $b_j(i, k)$  for all  $j \in \{1, 2, \dots, 48\}$ ,  $i \in \{1, 2, \dots, 48\}$  and  $k \in \{924, 925, 926\}$ .

For  $i \in \{1, 2, 4, 5, 13, 15, 19, 20, 21, 23, 24, 26, 27, 32, 33, 34, 38, 39, 41, 42, 43, 44, 45, 47\}$ , set  $c(i, 924) = 1$ ;

for other  $i$ , we run primality test and set  $c(i, 924) = 1$  for  $i \in \{8, 9, 16, 29, 30, 37, 40, 48\}$ .

For  $i \in \{3, 6, 7, 10, 11, 12, 14, 17, 18, 22, 25, 28, 31, 35, 36, 46\}$ , set  $c(i, 924) = 0$ .

For  $i \in \{1, 5, 6, 7, 8, 9, 10, 12, 14, 15, 18, 20, 21, 22, 26, 29, 30, 31, 33, 34, 35, 38, 39, 40, 42, 43, 45, 47\}$ , set  $c(i, 925) = 1$ ;

for other  $i$ , we run primality test and set  $c(i, 925) = 1$  for  $i \in \{11, 16, 19, 24, 25, 32, 37, 46, 48\}$ .

For  $i \in \{2, 3, 4, 13, 17, 23, 27, 28, 36, 41, 44\}$ , set  $c(i, 925) = 0$ .

For  $i \in \{2, 6, 7, 8, 9, 10, 12, 13, 16, 17, 21, 22, 23, 24, 26, 28, 30, 31, 32, 34, 35, 36, 37, 38, 39, 40, 41, 43, 47\}$ , set  $c(i, 926) = 1$ ;

for other  $i$ , we run primality test and set  $c(i, 926) = 1$  for  $i \in \{3, 18, 20, 45\}$ .

For  $i \in \{1, 4, 5, 11, 14, 15, 19, 25, 27, 29, 33, 42, 44, 46, 48\}$ , set  $c(i, 926) = 0$ .

Step 3: Generate a prime binary matrix  $P = (p(i, k))_{48 \times 3}$  based on  $C$ .

For  $i \in \{1, 2, \dots, 48\}$  and  $k \in \{924, 925, 926\}$ , set  $p(i, k) = 1 - c(i, k)$ .

Step 4: The prime numbers within the interval  $[11+210 \times 924, 211+210 \times 926]$  can be listed as

$$\begin{aligned} & \{ \alpha = r_i + 210k \mid p(i, k) = 1, i = 1, 2, \dots, 48, k = 924, 925, 926 \} \\ & = \{ 194471, 194263, 194057, 194267, 194269, 194479, 194483, 194069, 194071, \\ & \quad 194083, 194087, 194507, 194093, 194309, 194101, 194521, 194527, 194113, \\ & \quad 194323, 194119, 194543, 194141, 194353, 194149, 194569, 194371, 194581, \\ & \quad 194167, 194377, 194591, 194179, 194609, 194197, 194203, 194413, 194431, \\ & \quad 194647, 194443, 194653, 194239, 194659, 194671 \} \end{aligned}$$

### Appendix B: Theoretical Analysis and Computer Validation of Algorithm 2

The formulation of the Composite Binary Matrix  $C$  is the most vital step in our proposed prime listing algorithm, and Step 1 in Algorithm 2 locates a good portion of composites in a pre-specified interval with constant time complexity. In the following analysis, we will conduct a comprehensive computation to see the expectation of the composites located in a column by Step 1. It is more reasonable to count all composites within the period than to analyze the performance of each column, as the latter is too time-consuming to study.

After conducting Step 1, the entries of column  $c(i, k^*)$  are either 1 or 0. If  $c(i, k_1) - c(i, k_2) = 0$  for  $i \in \{1, 2, \dots, 48\}$ , we call the “result” of columns  $k_1$  and  $k_2$  identical. To analyze the effect of Step 1, we first elaborate that the “results” of Step 1 appear cyclically in the columns.

Theorem 5 implies that in a certain column  $k$  of matrix  $C$ , at which the entries are located as composite by Step 1 is determined by the result of

$$w_1 = k - r_1 \frac{k^*}{r_1}, w_2 = k - r_2 \frac{k^*}{r_2}, \dots, w_{48} = k - r_{48} \frac{k^*}{r_{48}}, \tag{11}$$

which is identical to

$$k \equiv w_j \pmod{r_j}, w_j \in \{0, 1, \dots, r_j - 1\}, j \in \{1, 2, \dots, 48\} \tag{12}$$

in modular arithmetic. If there exists  $j$  such that  $b_j^0(i, w_j) = 1$ , then  $b_j(i, k) = 1$ ,  $c(i, k) = 1$  and  $r_i + 210 \times k$  is a composite number.

We claim that there are infinite  $k$ 's satisfying Equation (11) and they are of the same cyclic form. Denoting  $L$  as the least common multiple of  $r_1, r_2, \dots, r_{48}$ , we have  $L = 1, 120, 700, 227, 446, 262, 311, 648, 514, 688, 739, 473, 686, 196, 699, 653, 055, 999, 250, 791, 866, 994, 591, 464, 755, 146, 178, 654, 953$ . Given  $w_1, w_2, \dots, w_{48}$ , the Chinese Remainder Theorem says that there's a unique solution  $k^*$ , if any, in  $[1, L]_N$  for Equation (11) and the general solution can be expressed as  $k^* + L \times t$ , where  $t \in N_0$ .

With the cyclic behavior detected, we can now limit our attention to the case where  $k \in [1, L]_N$ . However, the period  $L$  is too large to be counted directly, so we need some further analysis.

We should notice that for any  $i \in \{1, 2, \dots, 48\}$ , there's one and only one  $w_j$  in  $\{0, 1, \dots, r_j - 1\}$  such that  $b_j^0(i, w_j) = 1$ , indicating  $r_i + 210 \times k$  is composite. We denote this  $w_j$  as  $d_{ij}$  in the following discussion.

To avoid duplicated computation induced by the 5 composite numbers in  $S$ , we divide the set  $S$  into  $S_1 = \{11, 13, 143, 187, 209\}$ ,  $S_2 = \{17, 19, 121, 169\}$  and  $S_3 = S \setminus (S_1 \cup S_2)$ . Accordingly, the index set  $I = \{1, 2, \dots, 48\}$  is divided into  $I_1 = \{1, 2, 32, 42, 47\}$ ,  $I_2 = \{3, 4, 27, 38\}$  and  $I_3 = I \setminus (I_1 \cup I_2)$ .

The following Corollary demonstrates that  $\{w_j \mid j \in I_2 \cup I_3\}$  can uniquely determine a solution in  $[1, L]_N$  for Equation (11). The reason is that  $\{w_j \mid j \in I_1\}$  can be directly calculated given  $\{w_j \mid j \in I_2\}$ .

**Corollary of the Chinese Remainder Theorem.** Given  $w_1, w_2, \dots, w_{48}$ , Equation (11) has a solution in  $[1, L]_N$  if and only if

$$\begin{cases} w_1 \equiv w_{27} \pmod{11} \\ w_2 \equiv w_{38} \pmod{13} \\ w_{32} \equiv 78w_1 + 66w_2 \pmod{143} \\ w_{42} \equiv 34w_1 + 154w_3 \pmod{187} \\ w_{47} \equiv 133w_1 + 77w_4 \pmod{209} \end{cases} \tag{13}$$

and

$$k \equiv w_j \pmod{r_j}, w_j \in \{0, 1, \dots, r_j - 1\}, j \in I_2 \cup I_3 \tag{14}$$

has a solution in  $[1, L]_N$ .

The above Corollary is derived from the Chinese Remainder Theorem and can be easily verified. With this theorem, we can count how many composites are located by Step 1 in columns 1 to  $L$ .

Without loss of generality, we first focus on  $i = 1$ .

Since each column in  $[1, L]_N$  can be uniquely determined by  $\{w_j \mid j \in I_2 \cup I_3\}$ , we first consider  $\{w_j \mid j \in I_2\} = \{w_3, w_4, w_{27}, w_{38}\}$ . Notice that  $\{w_j \mid j \in I_1\}$  is determined by Equation (13) for any fixed  $\{w_j \mid j \in I_2\}$ .

If there exists at least a  $j \in I_2$  such that  $w_j = d_{1j}$ , then  $r_i + 210 \times k$  is a composite number no matter the value of  $w_j, j \in I_3$ . The number of  $\{w_3, w_4, w_{27}, w_{38}\}$  satisfying this case is

$r_3 r_4 r_{27} r_{38} - (r_3 - 1)(r_4 - 1)(r_{27} - 1)(r_{38} - 1) = 798947$ . In this case,  $w_j, j \in I_3$  can be chosen arbitrarily from the domain. Therefore, there are  $\prod_{j \in I_3} r_j$  composites located for each such  $\{w_3, w_4, w_{27}, w_{38}\}$  in  $[1, L]_N$ .

If  $w_j \neq d_{1j}$  for any  $j \in I_2$  but there exists at least a  $j \in I_1$  such that  $w_j = d_{1j}$ , there are also  $\prod_{j \in I_3} r_j$  composites located for each such

$\{w_3, w_4, w_{27}, w_{38}\}$ . Here  $w_j, j \in I_1$  is computed by Equation (13). The number of  $\{w_3, w_4, w_{27}, w_{38}\}$  satisfying this case is 864000 (Counted by computer, this value is the same when  $i = 2, 3, \dots, 48$ ).

If  $w_j \neq d_{1j}$  for any  $j \in I_1 \cup I_2$ , then a composite number is located if and only

if there exists  $w_j = d_{1j}$  for some  $j \in I_3$ . Therefore, there are  $\prod_{j \in I_3} r_j - \prod_{j \in I_3} (r_j - 1)$

composites located in  $[1, L]_N$  for each such  $\{w_3, w_4, w_{27}, w_{38}\}$ . The number of

$\{w_3, w_4, w_{27}, w_{38}\}$  satisfying this case is 4942080.

When  $i = 2, 3, \dots, 48$ , the number of composites located by Step 1 in columns  $[1, L]_N$  is the same as  $i = 1$ . Since

$$\prod_{j \in I_3} r_j = 169,673,829,864,171,987,737,296,863,243,628,479,671,120,141,228, \\ 188,658,546,265,896,353,105,711,020,739$$

and

$$\prod_{j \in I_3} (r_j - 1) = 102,585,912,060,381,157,597,756,546,764,390,409,686,966,581, \\ 524,267,485,974,696,558,592,000,000,000,000,$$

the total number of composites located by Step 1 in columns  $[1, L]_N$  is

$$T = 48 \times \left[ 798947 \prod_{j \in I_3} r_j + 864000 \prod_{j \in I_3} r_j + 4942080 \left( \prod_{j \in I_3} r_j - \prod_{j \in I_3} (r_j - 1) \right) \right] \\ = 29,458,197,272,202,902,414,774,112,677,093,526,253,964,920, \\ 792,154,274,897,889,854,062,645,442,967,016,575,437,744$$

The average number of composites located by Step 1 in a PTP's column is  $T/L = 26.285528057158732$  in the first period  $[1, L]_N$  and can be regarded as the overall expectation of the number of composites located by Step 1 in a PTP's column, which is over  $26.28/48 = 54.75\%$  of the numbers in the column.

To further validate this result, the number of composites located by Step 1 in PTP's columns 1 to  $10^{12}$  is on average 26.285525173295, which is quite close to our theoretical result.

The theoretical analysis and numerical experiment above illustrate that the proposed method largely reduces the number of primality checks in the prime listing task over a pre-specified interval.